Robust Scheduling for Wireless Charger Networks

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Abstract—In this paper, we deal with the problem of Robust scheduling for wireless charger networks (RULE), i.e., given a number of rechargeable devices, each of which may drift within a certain range, and a number of directional chargers with fixed positions and adjustable orientations distributed on a 2D plane, determining the orientations of the wireless chargers to maximize the overall expected charging utility while taking the charging power jittering into consideration. To address the problem, we first model the charging power as a random variable, and apply area discretization technique to divide the charging area into several subareas to approximate the charging power as the same random variable in each subarea and bound the approximation error. Then, we discretize the orientations of chargers to deal with the unlimited searching space of orientations with performance bound. Finally, by proving the submodularity of the problem after the above transformations, we propose an algorithm that achieves $(2 - \varepsilon)$-approximation ratio. We conduct both simulation and field experiments, and the results show that our algorithm can perform better than other comparison algorithms by 103.25% on average.

I. INTRODUCTION

Wireless Power Transfer (WPT) technology demonstrates its importance in our daily life due to its convenience such as no wiring, no contact, reliable and continuous power supply, and ease of maintenance, and attracts attentions from not only academic research but also industrial field. Wireless Power Consortium [1], which aims to promote the standardization of WPT, has grown to include 606 companies in 2018, and this number is more than twice of that of last year. In a WPT system, chargers transfer power to rechargeable devices via wireless with reasonable efficiency. Almost all existing works regarding WPT systems focus on performance optimization issues in determined environments, such as deploying specified number of chargers to maximize charging utility/flow rate for pre-deployed sensor networks [2]–[5].

In practice, however, the charging environments are always highly dynamic with uncertainties. A WPT system should be sufficiently robust to deal with such dynamics. Such dynamics are mainly due to the following reasons. First, instead of the traditional static WSNs [6], [7], the wireless devices may drift since they are not absolutely fixed; example cases include but not limited to: (1) devices can drift in their task areas to expand the scope of monitoring [8]; (2) sensors in pipe networks may drift when flow passes [9]; (3) non-fixed sensors [10] like underwater sensors [11] and sensors on unmanned aerial vehicles [12] may drift without control; (4) the sensors on bridges to monitor structural health are required to detect acceleration, ambient vibration and so on [13] so they may slightly drift with vibration. These drift cases may cause these devices out of charging area and decrease the network lifetime if cases are in wireless rechargeable sensor networks. Second, the charging power jitters [14], which means the charging power is not a certain value for a certain pair of charger and device. Unfortunately, recent works about wireless charger scheduling always do not consider these uncertainties.

In this paper, we deal with the problem of Robust scheduling for wireless charger networks (RULE), i.e., given a number of rechargeable devices, each of which may drift within a certain range, and a number of directional chargers with fixed positions and adjustable orientations distributed on a 2D plane, determining the orientations of the wireless chargers to maximize the overall charging utility while taking the charging power jittering into consideration.

Related works about our problem mainly involve robust wireless charging problem and wireless charger placement/scheduling problem. The former work only considers the jittering of electromagnetic radiation (EMR) rather than charging power while the latter works just apply deterministic models, which are not fit to our problem.

There are four main challenges in our problem. First, the charging power jitters rather than be a static value for a point on the plane, which raises the challenge to evaluate the charging power and even the charging utility. Second, the charging power is nonlinear with distance and it is additive from different chargers for one device, thus, the problem cannot be regarded as a simple geometric coverage problem. Third, it is a continuous problem that we should take the whole area into consideration where these devices can drift rather than only consider limited number of positions of devices as in static topologies. Fourth, infinite orientations for chargers to choose leads to the unlimited solution space.

To address the problem, for the first challenge, we model the charging power as a random variable and these random variables of the charging power from different chargers are taken into consideration when we evaluate the charging power. The former work only considers the jittering of electromagnetic radiation (EMR) rather than charging power while the latter works just apply deterministic models, which are not fit to our problem.

In this paper, we deal with the problem of Robust scheduling for wireless charger networks (RULE), i.e., given a number of rechargeable devices, each of which may drift within a certain range, and a number of directional chargers with fixed positions and adjustable orientations distributed on a 2D plane, determining the orientations of the wireless chargers to maximize the overall charging utility while taking the charging power jittering into consideration.

To address the problem, for the first challenge, we model the charging power as a random variable and these random variables of the charging power from different chargers are independent. We take the expectation of charging utility, a monotone increasing concave function of charging power, as the measurement. For the second one, we apply area discretization to divide the charging area of chargers into several subareas to approximate the charging power as the same random variable in each subarea.
The drifting areas of devices can also be divided into several subareas based on area discretization due to geometric symmetry, where expected charging power and charging utility are constant. Thus, we only need to calculate the expected charging utility in each subarea. For the fourth one, we discretize the orientations of chargers into limited ones with charging utility in each subarea. For the fifth one, we assume that the charging angle \( \alpha \) of chargers can be anywhere in the disk with a uniform distribution due to the practical concerns.

We build our charging model based on empirical studies [2], [3]. The charging area of a directional charger can be modeled as a sector with radius \( D \). As shown in Figure 1, charger \( s_i \) has its orientation angle \( \varphi_i \) and a charging angle \( \alpha \). We consider the typical case that the charging angle \( \alpha \in (0, \frac{\pi}{2}) \) due to the strong directional property of directional chargers. For example, the TX91501 power transmitter produced by Powercast [26] provides a charging angle of about 60° [2]. If the position \( p = (x, y) \) of the device \( o_j \) is in the charger’s charging area, it can always receive non-zero charging power, and none for otherwise. In Figure 1, \( o_j \) can always harvest non-zero charging power at position \( p_j, 1 \) but cannot at position \( p_j, 2 \). In addition, the charging power from a charger jitters due to multipath effect, so that we model it as a random variable [14]. By combining the widely used directional charging model and the probabilistic power model in [14], we can formalize our charging model as follows.

\[
\begin{align*}
\mathcal{N}
\left( \frac{a_1}{d(s_i, p) + b_1}, \frac{a_2}{(d(s_i, p) + b_2)^2} \right),
\end{align*}
\]

where \( a_1, b_1, a_2, b_2 \) are four parameters decided by the hardware and the environment and \( d(s_i, p) \) denotes the distance between charger \( s_i \) and the device position \( p \). The model shows that if the device is in the charging sector area of the charger, the receiving power obeys a normal distribution of which the parameters are related to the distance between the charger and the device. Moreover, we assume that the charging power from multiple chargers is additive, that is,

\[
\begin{align*}
P_w(p) = \sum_{i=1}^{N} P_w(s_i, \varphi_i, p).
\end{align*}
\]
equals 0, otherwise. To normalize, we suppose that $U(x)$ converges to 1 or finally reaches 1 without increasing any more. Since charging power is a random variable, charging utility of charging power is also a random variable. Here we use a linear bounded function to define $U(.)$.

$$U(x) = \begin{cases} 0, & x \leq 0, \\ x/P_{th}, & 0 < x \leq P_{th}, \\ 1, & x > P_{th}, \end{cases} \tag{3}$$

where $P_{th}$ is the power threshold which denotes rechargeable devices cannot harvest power any more when they have already received $P_{th}$ power.

### C. Problem Formulation

First, according to Equation (2), the expected charging utility for a rechargeable device $o_j$ at point $p = (x, y) \in O_j$ is $\mathbb{E}[U(\sum_{i=1}^{N} P_w(s_i, \varphi_i, (x, y)))]$ since there are $N$ chargers in total. Then, the overall expected charging utility for device $o_j$ in its DDA $O_j$ is the normalized sum of the expected charging utility of all the points in $O_j$. Since there are infinite number of points in an area, it becomes a normalized double integral in area $O_j$, i.e.,

$$\frac{1}{\pi r^2} \int_{(x, y) \in O_j} \mathbb{E}[U(\sum_{i=1}^{N} P_w(s_i, \varphi_i, (x, y)))] dxdy.$$  

Finally, there are $M$ devices in total, so we add their overall expected charging utility in their own moving circles together and get the average value. Thus, the problem is formalized as:

$$(P1) \max_{\varphi_i} \frac{1}{M \pi r^2} \sum_{j=1}^{M} \int_{(x, y) \in O_j} \mathbb{E}[U(\sum_{i=1}^{N} P_w(s_i, \varphi_i, (x, y)))] dxdy,$$

s.t. $O_j \subseteq \Omega$, $\varphi_i \in [0, 2\pi)$, $i = 1, ..., N$, $j = 1, ..., M.$ \tag{4}

The hardness of the problem is described in the theorem below.

**Theorem III.1.** The problem $P1$ is NP-hard.

We omit some proofs in this paper due to the space limit.

### IV. Solution

In this section, we introduce the detailed solution of the problem $P1$ which achieves approximation ratio $\frac{1}{2} - \epsilon$. We first apply area discretization to approximate the charging power in the same subarea with performance guarantee to address the nonlinearity of the problem. Then, to confine the unlimited searching space, we also discretize the orientations and bound the covering area gap. Finally, we reformulate the problem as maximizing a monotone submodular function subject to a partition matroid, propose our RULE algorithm, and give the overall performance guarantee.

#### A. Area Discretization

1) Piecewise Constant Function Approximation: To deal with the continuous and nonlinear charging power expectation, we first approximate the charging power of a charger by a piecewise constant function. Let $P_w(d)$ denote the charging power from a charger to a device with distance $d$. The piecewise constant function is defined as follows:

$$\tilde{P}_w(d) = \begin{cases} P_w(l(1)), & d = 0, \\ P_w(l(k)), & l(k-1) < d \leq l(k), \ (k = 1, ..., K), \\ 0, & d > D, \end{cases} \tag{5}$$

where $l(K) = D$. As shown in Figure 2, the charging area of charger $s_i$ is divided into two subareas and the charging power is approximated as $P_w(l(k))$ at any point in subarea $k$ $(k = 1, 2)$, so that the expected charging power in each subarea is a constant as the horizontal line shows. We use $\tilde{P}_w(d) = 0$ to denote the fact that the random variable $\tilde{P}_w(d)$ can only equal to 0, i.e., $P[\tilde{P}_w(d) = 0] = 1$. It is obvious that $\tilde{P}_w(d) = 0$ if and only if $P_w(d) = 0$. In the following lemmas corresponding to approximation error, we only discuss the case that the charger provides non-zero charging power.

2) Charging Utility Approximation for a DDA: The DDA $O_j$ is divided into several subareas by the piecewise constant function in which the approximated expectation and variance of charging power are constant. We denote the $k$th subarea in $O_j$ as $O_j^k$.

**Lemma IV.1.** Let $\tilde{P}_w(s_i, \varphi_i, O_j^k)$ denote the approximated charging power received at any point $p = (x, y) \in O_j^k$ from charger $s_i$ with orientation $\varphi_i$. By setting $l(0) = 0$,

$$l(k) = \min\{\sqrt{1 + \epsilon_1} (l(k-1) + b_1), \sqrt{1 + \epsilon_1} (l(k-1) + b_2)\},$$

where $k = 1, ..., K - 1$ and $l(K) = D$ such that

$$l(K-1) < D \leq \min\{\sqrt{1 + \epsilon_1} (l(K-1)+b_1), \sqrt{1 + \epsilon_1} (l(K-1)+b_2)\},$$

the approximation error of the expectation of charging utility from all the chargers can be bounded as:

$$1 \leq \frac{\mathbb{E}[U(\sum_{i=1}^{N} P_w(s_i, \varphi_i, p))]}{\mathbb{E}[U(\sum_{i=1}^{N} \tilde{P}_w(s_i, \varphi_i, O_j^k))]} \leq 1 + \epsilon_1. \tag{6}$$

3) Problem Reformulation: After charging utility approximation, the problem $P1$ is approximated as:

$$(P2) \max_{\varphi_i} \frac{1}{M \pi r^2} \sum_{j=1}^{M} \sum_{O_j^k \subseteq O_j} \mathbb{E}[U(\sum_{i=1}^{N} \tilde{P}_w(s_i, \varphi_i, O_j^k))] A(O_j^k),$$

s.t. $O_j \subseteq \Omega$, $\varphi_i \in [0, 2\pi)$, $i = 1, ..., N$, \tag{7}

where $A(O_j^k)$ denotes the area of $O_j^k$ and $\sum_{O_j^k \subseteq O_j} A(O_j^k) = \pi r^2$. $\tilde{P}_w(s_i, \varphi_i, O_j^k)$ denotes the approximated charging power at any point in subarea $O_j^k$ provided by charger $s_i$ with orientation $\varphi_i$.

### B. Orientation Discretization

Since the orientation of a charger is a continuous value and the relationship between all the covering subareas and the orientation is difficult to describe, we discretize the orientation into limited number of directions with $\Delta \varphi$ interval and bound the approximation error. We assume $\Delta \varphi$ can be divided by $2\pi$ with no remainder for simplicity.
Lemma IV.2. The gap between the original area and the area after orientation discretization covered by one charger in a DDA is bounded as $2r_o(D + r_o)\Delta \varphi$ when $\alpha \in [0, \frac{\pi}{2}]$.

Proof: We prove this lemma under polar coordinate system. First, we give notations and equations of lines and the DDA circle bound. Then, we discuss the three conditions: the charger outside/on/inside the DDA circle bound, respectively. At last, we derive the result.

Since we only consider a pair of charger and device, we do not care the indices of the charger and the device, and write $A(s_i, \varphi_j, O_j)$ as $A(\varphi)$ for simplicity. Suppose the charger is at the pole $(0,0)$ and the device position is $(r_o, 0)$. It should be noticed that any other cases that the device center position is not on the polar axis can be transformed into this situation by rotating the polar axis. The two beams of the charging area are $\theta = \varphi - \frac{\alpha}{2}$ and $\theta = \varphi + \frac{\alpha}{2}$. The DDA circle bound can be described as follows due to cosine theorem:

$$\rho^2 + \rho^2 - 2\rho \rho_o \cos \theta = r_o^2.$$  
(8)

We need to get the function $\rho(\theta)$ to further calculate the covering area. Since the circle is a quadratic equation about $\rho$, the equation may have two, one, or none root(s), which corresponds to different calculating cases.

We need to explore the maximum change rate of charging area while rotating the charger, that is, the maximum value of $|A'(\varphi)|$. It is obvious that the absolute value of the minimum and maximum $A'(\varphi)$ are the same due to the symmetry of the circle. In most cases, we can calculate the maximum value of $|A'(\varphi)|$ with the $\varphi$ which satisfies $A''(\varphi) = 0$.

We only consider the two line bounds of the charging area and neglect the farthest arc bound while rotating, since the farthest arc bound of charging area will contribute nothing to the area changing rate so that the area changing rate is smaller.

Charger outside the DDA: In this case, three situations may occur:
1) Both $\theta = \varphi - \frac{\alpha}{2}$ and $\theta = \varphi + \frac{\alpha}{2}$ cross the circle;
2) Only one beam $\theta = \varphi - \frac{\alpha}{2}$ or $\theta = \varphi + \frac{\alpha}{2}$ cross the circle;
3) None of the two beams cross the circle.

These situations can be seen in Figure 3(a)-(c), respectively. In situation 3), the covering area is either the area of the whole circle or 0, and both the changing rates are 0. We only consider the first and the second situations.

Define the two functions of the circle as

$$\rho_1(\theta) = \rho_o \cos \theta + \sqrt{r_o^2 - \rho_o^2 \sin^2 \theta},$$  
(9)

$$\rho_2(\theta) = \rho_o \cos \theta - \sqrt{r_o^2 - \rho_o^2 \sin^2 \theta},$$  
(10)

which are derived from Equation (8). In situation 1), the covering area can be calculated as:

$$A(\varphi) = \int_{\varphi - \frac{\alpha}{2}}^{\varphi + \frac{\alpha}{2}} \frac{1}{2} \rho_o^2(\theta)^2 d\theta$$  
(11)

$$= 2\rho_o \int_{\varphi - \frac{\alpha}{2}}^{\varphi + \frac{\alpha}{2}} \cos \theta \cdot \sqrt{r_o^2 - \rho_o^2 \sin^2 \theta} d\theta,$$

where $-\arcsin \left(\frac{\rho_o}{r_o}\right) < \varphi - \frac{\alpha}{2} < \varphi + \frac{\alpha}{2} < \arcsin \left(\frac{\rho_o}{r_o}\right)$.

Let $A''(\varphi) = 0$ and we get the only case $\sin(\varphi - \frac{\alpha}{2}) = \sin(\varphi + \frac{\alpha}{2})$, but the value of $\varphi$ exceeds its range. Thus, $A''(\varphi)$ must be monotone in its domain. Due to the symmetry, the values of $|A'(\varphi)|$ at the end points of the domain, i.e., $\varphi = -\arcsin \left(\frac{\rho_o}{r_o}\right) + \frac{\alpha}{2}$ and $\varphi = \arcsin \left(\frac{\rho_o}{r_o}\right) - \frac{\alpha}{2}$, are the same and this value is the upper bound in this situation. That is,

$$|A'(\varphi)| < 2\rho_o \cos(\alpha - \arcsin \left(\frac{\rho_o}{r_o}\right)) \sqrt{r_o^2 - \rho_o^2 \sin^2 (\alpha - \arcsin \left(\frac{\rho_o}{r_o}\right))}$$

$$\leq 2\rho_o r_o.$$

Charger on the DDA circle bound: In this case, the function of DDA circle bound reduces to:

$$\rho(\theta) = 2r_o \cos \theta, \quad \theta \in \left[-\frac{\pi}{2}, \frac{\pi}{2}\right],$$  
(13)

and it is easy to get the bound $|A'(\varphi)| \leq 2\rho_o r_o$. We omit the proof of this case here.

Charger inside the DDA: The covering case can be seen in Figure 5. In this case, we see that there must be one root for $\rho$ in Equation (8) since $r_o > \rho_o$ and we can write $\rho(\theta)$ as:

$$\rho(\theta) = \rho_o \cos \theta + \sqrt{\rho_o^2 \cos^2 \theta + (r_o^2 - \rho_o^2)}.$$  
(14)

The covering area $A(\varphi)$ can be calculated as:

$$A(\varphi) = \int_{\varphi - \frac{\alpha}{2}}^{\varphi + \frac{\alpha}{2}} \frac{1}{2} \rho^2(\theta)^2 d\theta$$  
(15)

$$= \frac{1}{2} \int_{\varphi - \frac{\alpha}{2}}^{\varphi + \frac{\alpha}{2}} \left(\rho_o \cos \theta + \sqrt{\rho_o^2 \cos^2 \theta + (r_o^2 - \rho_o^2)}\right)^2 d\theta.$$
Algorithm 1: Robust Scheduling

**Input:** Candidate orientation set of the $i$th charger $X_i$ ($i = 1, ..., N$), charger set $S = \{s_1, ..., s_N\}$, device set $O = \{o_1, ..., o_M\}$, density function of charging power provided by one charger, utility function $U(x)$, objective function $f(X)$.

**Output:** Selected orientation set $\Gamma$.

1. $\Gamma = \emptyset$
2. $X = \bigcup_{i=1}^{N} X_i$
3. while $|\Gamma| < N$ do
   4. $e^* = \arg \max_{e \in X \setminus \Gamma} f(X \cup \{e\}) - f(X)$
   5. $\Gamma = \Gamma \cup \{e^*\}$
   6. $X = X \setminus X_i$ where $e^* \in X_i$.

Define

$$A_f(x) = \rho_0 x \left( \rho_0 x + \sqrt{\rho_0^2 x^2 + (r_o^2 - \rho_0^2)} \right),$$

$$x \in [-1, 1], 0 < \rho_0 < r_o.$$  \hfill (16)

Then,

$$A'(\phi) = A_f(\cos(\phi + \frac{\alpha}{2})) - A_f(\cos(\phi - \frac{\alpha}{2}))$$

$$\leq A_f(1) - A_f(-1) = 2\rho_0 r_o,$$

since $A_f(x) = 0$ requires $(\rho_0^2 - r_o^2)^2 = 0$ which cannot be satisfied, and $A_f'(1) > 0, A_f'(x) > 0$ when $x \in [-1, 1]$. Thus, $A_f(x) (x \in [-1, 1])$ is a monotone increasing function, so we have the upper bound $2\rho_0 r_o$.

Therefore, in all the three cases, the area changing rate is no larger than $2\rho_0 r_o$. Considering that a charger covering a DDA with non-zero area requires $\rho_0 < D + r_o$, the area changing rate is less than $2r_o(D + r_o)$. Thus, the maximum changing area is less than $2r_o(D + r_o)\Delta \phi$ by rotating $\Delta \phi$. \hfill \blacksquare

After orientation discretization, the problem P2 in Equation (7) is reformulated as:

$$\max_{\phi_i} \frac{1}{M\pi r_o^2} \sum_{j=1}^{M} \sum_{O_j \subseteq O} \mathbb{E}[U(\sum_{i=1}^{N} \tilde{P}_w(s_i, \tilde{\phi}_i, O_j^i))]A(O_j^i),$$

subject to $O_j \subseteq \Omega$, $\tilde{\phi}_i = n_i \Delta \phi$,

$$n_i \in \{0, ..., \frac{2\pi}{\Delta \phi} - 1\}, \quad i = 1, ..., N, \quad j = 1, ..., M.$$  \hfill (18)

C. Algorithm and Theoretical Analysis

In this subsection, we will give a $(\frac{1}{2} - \epsilon)$-approximation algorithm and its further analysis. The details of the algorithm is shown in Algorithm 1, which is essentially a greedy algorithm. In the following, we will prove the validity of this algorithm by giving the approximation ratio.

**Definition IV.1.** [29] Let $S$ be a finite ground set. A real-valued set function $f : 2^S \rightarrow \mathbb{R}$ is normalized, monotonic and submodular if and only if it satisfies the following conditions, respectively: (1) $f(\emptyset) = 0$; (2) $f(S_1 \cup \{e\}) - f(S_1) \geq 0$ for any $S_1 \subseteq S$ and $e \in S \setminus S_1$; (3) $f(S_1 \cup \{e\}) - f(S_1) \geq f(S_2 \cup \{e\}) - f(S_2)$ for any $S_1 \subseteq S_2 \subseteq S$ and $e \in S \setminus S_2$.

**Definition IV.2.** [29] Partition matroid: Given $S = \bigcup_{i=1}^{k} S_i$ is the disjiont union of $k$ sets, $l_1, l_2, ..., l_k$ are positive integers, a partition matroid $M = (S, I)$ is a matroid where $I = \{X \subset S : |X \cap S_i| \leq l_i \text{ for } i \in [k]\}$. Based on the definitions above, problem P3 can be rewritten as:

$$\max f(X) = \frac{1}{M\pi r_o^2} \sum_{j=1}^{M} \sum_{O_j \subseteq O} \mathbb{E}[U(\sum_{\phi_i \in X} \tilde{P}_w(s_i, \tilde{\phi}_i, O_j^i))]A(O_j^i),$$

subject to $O_j \subseteq \Omega$, $X \in L, L = \{X \subseteq \Gamma : |X \cap \Gamma_q| \leq 1 \text{ for } q \in [N]\}$. \hfill (19)

**Lemma IV.3.** The objective function $f(X)$ in problem P4 is a monotone submodular function, and the constraint is a partition matroid constraint.

**Proof:** To verify the submodularity of function $f(X)$, we only need to check whether $f(X)$ satisfies the three conditions in Definition IV.1. Define

$$f_1(X, j, k) = \mathbb{E}[U(\sum_{\tilde{\phi}_i \in X} \tilde{P}_w(s_i, \tilde{\phi}_i, O_j^i))].$$

Thus, we have

$$f(X) = \frac{1}{M\pi r_o^2} \sum_{j=1}^{M} \sum_{O_j \subseteq O} f_1(X, j, k)A(O_j^i).$$  \hfill (21)

Note that the area discretization is different for different charger set $X$, i.e., $O_j^i$'s are different. To unify the subareas, we discretize the subareas more finely to ensure the different charger sets in the same condition give common subareas.

Further, we define the probability density function for random variable $\sum_{\tilde{\phi}_i \in X} \tilde{P}_w(s_i, \tilde{\phi}_i, O_j^i)$ as $g_X(x)$. Since
charging power provided by different chargers is independent, the random variables $\sum_{s_1} P_{w}(s_1, \hat{\varphi}, O_k^j)$ and $\sum_{s_2} P_{w}(s_2, \hat{\varphi}, O_k^j)$ are independent if $S_1 \cap S_2 = \emptyset$. That is, the joint probability density function $\tilde{g}_{S_1 \cup S_2}(x, y)$ of the random variable $\sum_{s_1} P_{w}(s_1, \hat{\varphi}, O_k^j) + \sum_{s_2} P_{w}(s_2, \hat{\varphi}, O_k^j)$ is equal to $\tilde{g}_{S_1}(x)\tilde{g}_{S_2}(y)$.

It is obvious that $f(X)$ satisfies the first and the second conditions with $f(\emptyset) = 0$ and monotonically increasing property. For the third condition, we give a lemma first.

**Lemma IV.4.** The monotonically increasing concave function $U(x)$ has the property that

$$F_{u}(x, y, z) = (U(x + z) - U(x)) - (U(x + y + z) - U(x + y)),$$

where $x, y, z \geq 0$.

By referring to the result in Lemma IV.4, we define

$$\Delta f_1(S_1, S_2, \{e\}, j, k) = (f_1(S_1 \cup \{e\}, j, k) - f_1(S_1, j, k)) - (f_1(S_2 \cup \{e\}, j, k) - f_1(S_2, j, k))$$

$$= \left[ E[U(\sum_{\hat{\varphi} \in S_1} P_{w}(s_1, \hat{\varphi}, O_k^j))] - E[U(\sum_{\hat{\varphi} \in S_1 \cup \{x\}} P_{w}(s_1, \hat{\varphi}, O_k^j))] \right]$$

$$- \left[ E[U(\sum_{\hat{\varphi} \in S_1 \cup \{y\}} P_{w}(s_1, \hat{\varphi}, O_k^j))] - E[U(\sum_{\hat{\varphi} \in S_1 \cup \{z\}} P_{w}(s_1, \hat{\varphi}, O_k^j))] \right]$$

$$= \int_{-\infty}^{+\infty} \int_{-\infty}^{+\infty} U(x + y + z)\tilde{g}_{S_1}(x)\tilde{g}_{S_2}(y)\tilde{g}_{S_1}(z)dzdydz - \int_{-\infty}^{+\infty} U(x)\tilde{g}_{S_1}(x)dx$$

$$= \int_{0}^{+\infty} \int_{0}^{+\infty} F_{u}(x, y, z)\tilde{g}_{S_1}(x)\tilde{g}_{S_2}(y)\tilde{g}_{S_1}(z)dzdydz \geq 0.$$  

Thus,

$$f(S_1 \cup \{e\}) - f(S_1) - (f(S_2 \cup \{e\}) - f(S_2))$$

$$\geq \frac{1}{M\pi r_{D}^2} \sum_{j=1}^{M} \sum_{O_k^j \subseteq O_j} \Delta f_1(S_1, S_2, \{e\}, j, k) A(O_k^j).$$  

(24)

That is,

$$f(S_1 \cup \{e\}) - f(S_1) \geq f(S_2 \cup \{e\}) - f(S_2), \quad S_1 \subseteq S_2 \subseteq S, \quad e \in S \backslash S_2.$$  

(25)

To sum up, $f(X)$ in Equation (19) is a monotone submodular function and clearly the constraint is a partition matroid constraint.
First, if the orientations in $\Gamma'_3$ are far from those in $\Gamma'_2$, there is certainly a solution $\Gamma'_4$ which is no better than $\Gamma'_3$, and the orientations in which are nearest to those in $\Gamma'_2$. Since $U(.)$ finally reaches 1, \[ E[U(\sum_{\varphi_j \in \Gamma'_2} \tilde{P}_w(s_i, \varphi_i, O'_j))] \leq 1. \]

According to Lemma IV.2, we know that the maximum changing of area of all the subareas covered by one charger is $2r_o(D + r_o) \Delta \varphi$ in one DDA. Thus, the maximum changing area in total for all $N$ chargers is $2Nr_o(D + r_o) \Delta \varphi$, i.e.,

\[ \sum_{O'_j \in O_j} (A(O'_j) - A(O'_j)) \leq 2Nr_o(D + r_o) \Delta \varphi \quad (30) \]

in every DDA. So, the result in Equation (29) stands.

We consider a real scenario that there are enough chargers to give the common charging area full charging utility, and the charger should be placed at the position where at least one center of a DDA is within distance $D$. The worst case is that all the chargers are placed at one point and with the same orientation, thus, makes the covering area harvest full charging utility. In this case, the average charging utility is $A_w/M\pi r_d^2$, where $A_w$ is the covering area as shown in Figure 6(a)-(d) of Cases 1-4, and it is calculated by Equation (26). The critical conditions to differentiate Case 2, Case 3, and Case 4 are shown in Figure 6(e). In Case 3, there are two intersections on each line bound of the charging sector and the minimum $\frac{4r_o}{\pi}$ is $\beta_1$ as shown in 6(e), where $\cos \beta_1 = \frac{D^2 + D^2 - r_o^2}{2xDxD} = 1 - \frac{r_o^2}{D^2}x$. so $\sin \frac{\beta_1}{2} = \sqrt{1 - \cos^2 \beta_1} = \frac{r_o}{\pi} \sqrt{1 - \frac{r_o^2}{D^2}}$. The maximum $\frac{4r_o}{\pi}$ is obviously $\beta_2$ as shown in Figure 6(e) where $\sin \beta_2 = \frac{r_o}{\pi} \sqrt{1 - \frac{r_o^2}{2D^2}}$, that two line bounds of the charging sector are tangent to the DDA. Thus, in Case 3, $\sin \beta_1 \leq \sin \frac{4r_o}{\pi} < \sin \beta_2$, that is, $\frac{4r_o}{\pi} \sqrt{1 - \frac{r_o^2}{D^2}} \leq \sin \beta_1 \leq \frac{4r_o}{\pi} \sqrt{1 - \frac{r_o^2}{2D^2}}$. It is obvious that

\[ \frac{1}{M\pi r_d^2} \sum_{j=1}^{M} \int_{(x,y) \in O_j} E[U(\sum_{\varphi_i \in \Gamma'_1} \tilde{P}_w(s_i, \varphi_i, (x,y)))] dx dy \geq \frac{A_w}{M\pi r_d^2}. \quad (31) \]

Thus,

\[ \frac{1}{M\pi r_d^2} \sum_{j=1}^{M} \int_{(x,y) \in O_j} E[U(\sum_{\varphi_i \in \Gamma'_1} \tilde{P}_w(s_i, \varphi_i, (x,y)))] dx dy \geq \frac{A_w}{M\pi r_d^2}. \]

From Equation (27), (28) and (32), we have

\[ \frac{1}{1 + \epsilon_1} \frac{1}{M\pi r_d^2} \sum_{j=1}^{M} \int_{(x,y) \in O_j} E[U(\sum_{\varphi_i \in \Gamma'_1} \tilde{P}_w(s_i, \varphi_i, (x,y)))] dx dy \]

\[ \leq \frac{1}{2^{2}} \frac{1}{M\pi r_d^2} \sum_{j=1}^{M} \int_{(x,y) \in O_j} E[U(\sum_{\varphi_i \in \Gamma'_1} \tilde{P}_w(s_i, \varphi_i, O'_j))] dx dy \]

\[ \sum_{O'_j \in O_j} (A(O'_j) - A(O'_j)) \leq 2Nr_o(D + r_o) \Delta \varphi \]

\[ \leq \frac{1}{M\pi r_d^2} \sum_{j=1}^{M} \int_{(x,y) \in O_j} E[U(\sum_{\varphi_i \in \Gamma'_1} \tilde{P}_w(s_i, \varphi_i, O'_j))] dx dy. \]

Thus, we can bound the gap between the solutions to problem P4 and P1 as

\[ \frac{1}{M\pi r_d^2} \sum_{j=1}^{M} \sum_{\varphi_i \in \Gamma'_1} \tilde{P}_w(s_i, \varphi_i, O'_j)] A(O'_j) \]

\[ \leq \frac{1}{2} \left( \frac{1}{1 + \epsilon_1} - \frac{2M Nr_o(D + r_o)}{A_w} \Delta \varphi \right) \]

We omit the time complexity analysis to save space. $\blacksquare$

V. SIMULATION RESULTS

A. Evaluation Setup

In the simulations, we deploy both devices and chargers randomly in a 15 m x 15 m square area such that the devices can only move in this area, i.e., the moving centers of DDAs are randomly deployed in the center square with side length $(15 - 2r_o)$ m, and there must be at least one DDA center within
distance $D$ with respect to a charger. The default settings of the parameters are: $D = 6 \text{ m}$, $r_o = 2 \text{ m}$, $M = 8$, $N = 10$, $a_1 = 100$, $b_1 = 5$, $a_2 = 4$, $b_2 = 3$, $\epsilon = 0.1$, and $P_{th} = 1.5W$.

As there is no existing algorithm that addresses our problem, we propose three algorithms for comparison. Randomized Orientations (RO) generates the orientations of the chargers randomly. No Moving or Jittering (NMJ) generates the orientations of the chargers by DCS extraction for point case algorithm as proposed in [3] with the assumption that $r_o = 0$, i.e., the devices are fixed at their DDA centers, and there is no power jittering. Nearest Facing Device (NFD) selects the charger’s nearest DDA center and makes the charger face it.

All results in the figures are the average utility expectations of 100 random deployments. We use Utility Expectation or $\mathbb{E}[U]$ in the figures to denote the expectation of charging utility.

**B. Performance Comparison**

1) Impact of Number of Chargers $N$: Our simulation results show that on average RULE outperforms NFD, NMJ, and RO by 12.53%, 19.77%, and 68.91%, respectively, in terms of $N$. As shown in Figure 7, we test the expected charging utility with $N$ from 6 to 13. All of these four algorithms achieve higher charging utility expectation with larger number of chargers. The increasing trend slows down when $N$ becomes larger, since the expected charging utility is approximating the maximum expected charging utility of the specific topology gradually.

2) Impact of Number of Devices $M$: Our simulation results show that on average RULE outperforms NFD, NMJ, and RO by 12.85%, 16.42%, and 59.32%, respectively, in terms of $M$. We evaluate the expectation of charging utility when $M$ varies from 8 to 15. The charging utility expectation of RULE, NFD, and NMJ decreases slowly with $M$ becomes larger, while the results of RO fluctuate, as shown in Figure 8. Since the region is relatively small, the overall area of DDAs for chargers to cover will not become much larger as $M$ increases and there will be more overlapped area of DDAs. The charging utility expectation of NFD decreases faster and finally reaches that of NMJ, while that of RO remains low.

3) Impact of Charging Angle $\alpha$: Our simulation results show that on average RULE outperforms NFD, NMJ, and RO by 19.56%, 24.10%, and 103.25%, respectively, in terms of $\alpha$. Charging angle $\alpha$ varies from $34^\circ$ to $90^\circ$ and the charging utility expectation of all the four algorithms increases fast when $\alpha$ increases as shown in Figure 9. Still, the increasing trend becomes slow when $\alpha$ becomes larger.

**Table II**

<table>
<thead>
<tr>
<th>Index</th>
<th>Position</th>
<th>RULE</th>
<th>NFD</th>
<th>NMJ</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>(51.31, 20.09)</td>
<td>51.6$^\circ$</td>
<td>48.5$^\circ$</td>
<td>78.5$^\circ$</td>
</tr>
<tr>
<td>2</td>
<td>(140.29, 22.94)</td>
<td>101.1$^\circ$</td>
<td>101.1$^\circ$</td>
<td>101.1$^\circ$</td>
</tr>
<tr>
<td>3</td>
<td>(182.53, 62.25)</td>
<td>145.0$^\circ$</td>
<td>145.0$^\circ$</td>
<td>145.0$^\circ$</td>
</tr>
<tr>
<td>4</td>
<td>(50.13, 287.46)</td>
<td>327.5$^\circ$</td>
<td>327.5$^\circ$</td>
<td>327.5$^\circ$</td>
</tr>
<tr>
<td>5</td>
<td>(148.82, 290.67)</td>
<td>266.8$^\circ$</td>
<td>266.8$^\circ$</td>
<td>266.8$^\circ$</td>
</tr>
<tr>
<td>6</td>
<td>(273.39, 292.30)</td>
<td>245.4$^\circ$</td>
<td>245.4$^\circ$</td>
<td>245.4$^\circ$</td>
</tr>
</tbody>
</table>

4) Impact of DDA Radius $r_o$: Our simulation results show that on average RULE outperforms NFD, NMJ, and RO by 15.98%, 20.88%, and 54.36%, respectively, in terms of $r_o$. As shown in Figure 10, $r_o$ increases from 1.5 to 5 and the charging utility expectation decreases when $r_o$ increases and RULE, NFD, and NMJ decrease faster than RO.

**VI. FIELD EXPERIMENTS**

**A. Testbed**

As shown in Figure 11, our testbed consists of six TX91501 power transmitters produced by Powercast [26], [30]–[34], six rechargeable sensor nodes, and an AP connecting to a laptop to report the collected data. All the chargers and the devices are deployed in a $300 \text{ cm} \times 300 \text{ cm}$ square area, and the positions of the DDA centers are (124.61, 102.84), (235.97, 122.47), (179.47, 167.46), (299.54, 218.44), (198.95, 221.39), and (145.26, 226.89), and the positions of chargers are shown in Table II. The DDA radius $r_o$ is $50 \text{ cm}$ and $P_{th} = 30 \text{ mW}$. To evaluate the charging utility expectation of a device in its DDA, we measured 17 points in each DDA, every $25 \text{ cm}$ one point from the DDA center then add four points on the circle bound. We waited for at least $20 \text{ s}$ at each point to collect enough data to calculate the expectation of charging utility.

**B. Experimental Results**

The scheduling strategies for RULE, NFD, and NMJ are shown in Table II and the visual ones are shown in Figure 12-14. From Figure 15, we can see that our algorithm RULE allows the charging utility expectation of each device roughly equal and relatively large, while that of the other two algorithms varies much for these devices. The CDF plot of all the data collection points is shown in Figure 16, and it shows that the line of RULE approaches 1 at the lowest speed, which indicates that the expectation charging utility of RULE is generally more than the other two.
VII. CONCLUSION

In this paper, we deal with the robust scheduling for wireless charger networks problem. The key novelty of this paper is on proposing the first scheme for robust wireless charger scheduling considering power jittering and device drifting. We establish the probabilistic model for charging power, apply area and orientation discretization, and conduct both simulation and field experiments. The key technical depth is to reformulate the objective function to a submodular function and bound the overall performance gap. Our experimental results show that our proposed method achieves good performance and can outperform comparison algorithms by 103.25% on average.

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