Efficient Adaptive Online Learning via Frequent Directions

Yuanyu Wan and Lijun Zhang, Member, IEEE

Abstract—By employing time-varying proximal functions, adaptive subgradient methods (ADAGRAD) have improved the regret bound and been widely used in online learning and optimization. However, ADAGRAD with full matrix proximal functions (ADA-FULL) cannot handle large-scale problems due to the impractical $O(d^3)$ time and $O(d^2)$ space complexities, though it has better performance when gradients are correlated. In this paper, we propose two efficient variants of ADA-FULL via a matrix sketching technique called frequent directions (FD). The first variant named as ADA-FD directly utilizes FD to maintain and manipulate low-rank matrices, which reduces the space and time complexities to $O(\tau d)$ and $O(\tau^2 d)$ respectively, where $d$ is the dimensionality and $\tau \ll d$ is the sketching size. The second variant named as ADA-FFD further adopts a doubling trick to accelerate FD used in ADA-FD, which reduces the average time complexity to $O(\tau d)$ while only doubles the space complexity of ADA-FD. Theoretical analysis reveals that the regret of ADA-FD and ADA-FFD is close to that of ADA-FULL as long as the outer product matrix of gradients is approximately low-rank. Experimental results demonstrate the efficiency and effectiveness of our algorithms.

Index Terms—Online learning, frequent directions, adaptive subgradient methods.

1 INTRODUCTION

Online learning refers to the process of answering a sequence of questions given answers to previous questions, which enjoys an attractive combination of computational efficiency and theoretical guarantees [2]. Recently, it has received ever-increasing attention due to the emergence of large-scale applications such as online classification [3], [4], [5], [6], [7], online advertisements recommendation [8], [9], [10], online metric learning [11], [12], [13], [14], online matrix factorization [15], [16], [17], [18], online optimization [19], [20], [21], [22], [23], and online anomaly detection [24], [25]. Adaptive subgradient methods (ADAGRAD), which employ proximal functions that adapt to the geometry of the data observed in earlier iterations, are popular for online learning and optimization [26]. According to the type of proximal functions, ADAGRAD can be divided into learning with full matrix proximal functions (ADA-FULL) and learning with diagonal matrix proximal functions (ADA-DIAG). In contrast to ADA-FULL, ADA-DIAG has been successfully applied to many large-scale applications because of its light computations and storage.

However, the diagonal matrix maintained in ADA-DIAG only contains limited information of gradient outer products. This shortcoming causes that the regret of ADA-DIAG is worse than that of ADA-FULL when the high-dimensional data are dense and have an approximately low-rank structure. In consideration of this dilemma, Duchi et al. [26] proposed an open question concerning whether we can efficiently use full matrices in the proximal functions. To solve this problem, Krummenacher et al. [27] utilized random projections to approximate ADA-FULL, and developed two methods, namely ADA-LR and RADAGRAD. Compared with ADA-FULL, ADA-LR has the same space complexity $O(d^2)$ and a slightly reduced time complexity $O(\tau d^2)$, where $d$ is the dimensionality and $\tau \ll d$ is the number of random projections. Due to the quadratic dependence on $d$, ADA-LR is impractical for high-dimensional data, though it is equipped with a formal regret bound. In contrast, the time and space complexities of RADAGRAD are linear in $d$, but it lacks theoretical guarantees owing to further approximations. In our recent work [28], we also utilized random projections to accelerate ADA-FULL, and proposed ADA-DP that has complexities linear in $d$. However, its theoretical guarantees are non-deterministic, and the regret bound contains additional terms that never vanish.

To tackle the above limitations, we develop two efficient variants of ADA-FULL which are also principled. The computational bottleneck of ADA-FULL is to store the outer product matrix of gradients and compute its square root in each round. We note that a similar bottleneck also exists in other second order methods such as online Newton step (ONS) [29], and recently matrix sketching techniques have been used to reduce the computational cost of ONS [30]. Motivated by previous work, we first propose to directly employ frequent directions [31] to approximate the outer product matrix of gradients in ADA-FULL. By utilizing the low-rank structure, our first efficient variant of ADA-FULL, named ADA-FD, reduces the space complexity from $O(d^2)$ to $O(\tau d)$ and the time complexity from $O(d^2)$ to $O(\tau^2 d)$, where $\tau$ is the sketching size, which implies both the space and time complexities of ADA-FD are linear in the dimensionality $d$. Then, we further propose a much faster variant of ADA-FULL, named ADA-FFD, by accelerating frequent directions used in ADA-FD, which only needs to double the space complexity. Assuming $\tau \ll \sqrt{d}$, ADA-FFD reduces the average time complexity to $O(\tau d)$ that is linear.
in both the dimensionality $d$ and the sketching size $\tau$.

Moreover, because the idea of ADAGRAD can be incorporated into either the primal-dual subgradient method [32] or the composite mirror descent method [33], we also develop two efficient methods for both ADA-FD and ADA-FFD according to the type of subgradient methods, and prove that the regret bounds of our methods are close to that of ADA-FULL. The slight differences in the regret bounds are caused by the approximation error of frequent directions or fast frequent directions, and vanish when the sketching size is larger than the rank of gradient outer products. More generally, when the outer product matrix of gradients is approximately low-rank, a small number of sketching updates: more randomized approximations, they presented a more s-

\begin{align}
\beta_{t+1} &= \arg\min_\beta \left\{ \eta \left( \frac{1}{t} \langle \tilde{g}_t, \beta \rangle + \eta \varphi(\beta) \right) + \frac{1}{t} \Psi_t(\beta) \right\} \\
&= -\eta H_t^{-1} \tilde{g}_t, \quad \text{if } \varphi = 0.
\end{align}

And the composite mirror descent method updates by

\begin{align}
\beta_{t+1} &= \arg\min_\beta \left\{ \eta \langle g_t, \beta \rangle + \eta \varphi(\beta) + B \psi_t(\beta, \beta_t) \right\} \\
&= \beta_t - \eta H_t^{-1} \tilde{g}_t, \quad \text{if } \varphi = 0.
\end{align}

Because the storage complexity of $G_t$ is $O(d^2)$ and the time complexity of finding its square root is $O(d^3)$, ADA-FULL is impractical for large-scale problems.

To reduce the complexities of ADA-FULL, Krummenacher et al. [27] proposed two methods based on the fast randomized singular value decomposition (SVD) [34] to approximate the proximal term $\psi_t(\beta)$. Let $\Pi \in \mathbb{R}^{d \times d}$ be the random matrix of the subsampled randomized Fourier transform. They first used the following steps to update $\beta_t$:

\begin{align}
G_t &= G_{t-1} + g_t g_t^T \\
\tilde{G}_t &= G_t \Pi^T \\
Q &= \text{QR-decomposition} \quad \text{QR-projection} \\
\beta_{t+1} &= \beta_{t} - \eta V(\Sigma_{1/2} + \delta_t) - 1 V^T g_t
\end{align}

and the resulting algorithm is referred to as ADA-LR. Note that the space and time complexities of ADA-LR are respectively $O(d^2)$ and $O(d \tau^2)$, which prevent ADA-LR from being practical for high-dimensional data. By introducing more randomized approximations, they presented a more scalable algorithm RADAGRAD that performs the following updates:

\begin{align}
\tilde{g}_t &= \Pi g_t, \quad \text{Random projection} \\
\tilde{G}_t &= \tilde{G}_{t-1} + \tilde{g}_t \tilde{g}_t^T \\
\tilde{Q} &= \text{qr-update}(Q_{t-1}, R_{t-1}, \tilde{g}_t, \tilde{R}_t) \\
\tilde{V} &= Q_t W, \quad \gamma_t = \eta (\tilde{g}_t - \tilde{V}^T \tilde{g}_t) \\
\beta_{t+1} &= \beta_{t} - \eta (\Sigma_{1/2} + \delta_t)^{-1} V^T \tilde{g}_t - \gamma_t
\end{align}

where qr_update means QR decomposition with the rank-one update. In this case, RADAGRAD reduces the space and time complexities to $O(\tau d^2)$ and $O(\tau d)$ respectively. Unfortunately, it is a heuristic method and lacks theoretical guarantees. When $f_t(\beta_t) = l(\beta_t^T x_t)$ where $x_t$ is a data vector independent from the learning algorithm, we recently proposed ADA-DP [28] based on random projections to attain theoretical guarantees and keep complexities linear in $d$. However, due to the randomness of random projections, the regret bound of ADA-DP is non-deterministic and is worse than ADA-FULL even when $\tau = d$.

2.2 Matrix Sketching and Frequent Directions

For any given matrix $A \in \mathbb{R}^{d \times d}$, the purpose of sketching is to generate a matrix $B \in \mathbb{R}^{\tau \times d}$, where $\tau \ll t$ is the sketching size such that $A \approx B$ or $A^T \approx B^T B$. There
are many matrix sketching techniques including frequent directions [31], [35], [36], random projection [37], [38], [39], and column selection [40]. Frequent directions [31] is a deterministic matrix sketching technique by extending the well-known algorithm for approximating item frequencies in streams [41]. Let \( B = [b_1, b_2, \ldots, b_τ] = 0_{τ \times d} \) where \( τ \ll \min\{t, d\} \). For any given matrix \( A = [a_1, a_2, \ldots, a_t]^T \), frequent directions processes each row of \( A \) as follows

\[
\begin{align*}
\mathbf{b}_r = \mathbf{a}_r, & \quad \mathbf{UΣV}^T = \mathbf{B} \quad \text{SVD} \\
\mathbf{B} = \sqrt{\Sigma^2 - \sigma I_τV^T} & \quad \text{where } \sigma = \Sigma^2_{ττ}
\end{align*}
\]

and products a sketch matrix \( B \). Because the time complexity of processing each row is \( O(τ^2d) \) that is dominated by computing the SVD of \( B \), the total time complexity of frequent directions is \( O(τ^2dt) \) which suffers quadratic dependence on the sketching size \( τ \). To further reduce it to \( O(τdt) \), Ghashami et al. [31] have proposed fast frequent directions (FFD) at the expense of doubling the space complexity. Let \( B = [b_1, b_2, \ldots, b_τ]^T = 0_{2τ \times d} \). Fast frequent directions processes each row of \( A \) as follows

Insert \( a_τ \) into a zero valued row of \( B \)

if \( B \) has no zero valued rows then

\[
\begin{align*}
\mathbf{UΣV}^T = \mathbf{B} \quad \text{SVD} \\
\mathbf{B} = \sqrt{\max(\Sigma^2 - \sigma I_τ, 0)V^T} & \quad \text{where } \sigma = \Sigma^2_{ττ}
\end{align*}
\]

end if

where the if statement is only triggered once every \( τ + 1 \) rows. Therefore, the number of computing the SVD of \( B \) is only \( t/(τ + 1) \) and the total time complexity of fast frequent directions is \( O(τ^2dt/(τ + 1)) = O(τdt) \).

Recently, the techniques of matrix sketching have been used by Luo et al. [30] to accelerate online Newton step [29]. They studied ONS that updates by

\[
\Lambda_t = α \mathbf{I}_d + \sum_{i=1}^{t} η_i \mathbf{g}_i \mathbf{g}_i^T \quad \text{and } \beta_{t+1} = β_t - \Lambda_t^{-1} \mathbf{g}_t
\]

where \( α > 0 \) and \( η_t > 0 \) are some parameters for general convex functions, and used matrix sketching techniques to construct a low-rank approximation of the second order information. Motivated by Luo et al. [30], our work employs frequent directions to calculate a low-rank approximation of the full matrix, which obviously reduces the storage and time complexities of ADA-FULL. Compared to the approximation algorithm used by Krummenacher et al. [27], frequent directions has two advantages which are deterministic theoretical properties and easy implementations.

### 3 Efficient Variants of ADA-FULL

In this section, we first introduce our two efficient variants of ADA-FULL and the corresponding theoretical results. Then, we compare our work with Krummenacher et al. [27] and Luo et al. [30]. To facilitate presentations, we consider the case \( φ = 0 \), and our methods can be extended to the general case \( φ \neq 0 \).

#### 3.1 Adaptive Online Learning via FD

Define \( C_t = [g_1, g_2, \ldots, g_t]^T \in \mathbb{R}^t \times d \), where each row is a gradient vector. The outer product matrix of gradients can be calculated as

\[
\mathbf{G}_t = \sum_{i=1}^{t} \mathbf{g}_i \mathbf{g}_i^T = \mathbf{C}_t^T \mathbf{C}_t.
\]

To accelerate the computation of \( \mathbf{H}_t^{-1} \), we take advantage of frequent directions to produce a low-rank approximation of \( \mathbf{C}_t \) denoted by \( \mathbf{S}_t = [s_1, s_2, \ldots, s_t]^T \in \mathbb{R}^τ \times d \). Then we can redefine \( \mathbf{H}_t \) in the proximal term as

\[
\mathbf{H}_t = δ \mathbf{I}_d + (\mathbf{S}_t^T \mathbf{S}_t)^{1/2}.
\]

Let the SVD of \( \mathbf{S}_t \) be \( \mathbf{UΣV}^T \) where \( \mathbf{U} \in \mathbb{R}^{τ \times τ} \), \( \Sigma \in \mathbb{R}^{τ \times τ} \) and \( \mathbf{V} \in \mathbb{R}^{d \times τ} \), then we have \( (\mathbf{S}_t^T \mathbf{S}_t)^{1/2} = \mathbf{VΣV}^T \). By Woodbury formula [42], we have

\[
\mathbf{H}_t^{-1} = (δ \mathbf{I}_d + \mathbf{VΣV}^T)^{-1} = \frac{1}{δ}(\mathbf{I}_d - \mathbf{V}(δ \mathbf{I}_τ + \Sigma)^{-1} \mathbf{ΣV}^T).
\]

With the above procedures, we propose our first efficient variant of ADA-FULL, named as adaptive online learning via frequent directions (ADA-FD). Then, we first consider incorporating ADA-FD into the primal-dual subgradient method. According to the update rules in (1), in the \( t \)-th round, our algorithm performs the following updates

\[
\begin{align*}
\mathbf{s}_{t+1}^{g} & = \mathbf{g}_t, \quad \mathbf{g}_t = \mathbf{g}_t - \mathbf{g}_t \\
\mathbf{UΣV}^T & = \mathbf{S}_{t-1}, \quad \sigma_t = \Sigma^2_{ττ} \quad \text{SVD} \\
Σ' & = \sqrt{Σ^2 - σ I_τ}, \quad \Sigma_t = Σ'V^T \\
β_{t+1} & = β_t - \eta_t (\mathbf{g}_t - \mathbf{V}(δ I_τ + Σ')^{-1} Σ'V^T \mathbf{g}_t)
\end{align*}
\]

The detailed procedures of ADA-FD for the primal-dual subgradient method are summarized in Algorithm 1, and this method is named as adaptive dual averaging via frequent directions.

Moreover, we incorporate ADA-FD into the composite mirror descent method. According to the update rules in (2), in the \( t \)-th round, our algorithm performs the following updates

\[
\begin{align*}
\mathbf{s}_{t+1}^{g} & = \mathbf{g}_t \\
\mathbf{UΣV}^T & = \mathbf{S}_{t-1}, \quad \sigma_t = \Sigma^2_{ττ} \quad \text{SVD} \\
Σ' & = \sqrt{Σ^2 - σ I_τ}, \quad \Sigma_t = Σ'V^T \\
β_{t+1} & = β_t - \eta_t (\mathbf{g}_t - \mathbf{V}(δ I_τ + Σ')^{-1} Σ'V^T \mathbf{g}_t)
\end{align*}
\]

The detailed procedures of ADA-FD for the composite mirror descent method are summarized in Algorithm 2, and this method is named as adaptive mirror descent via frequent directions.

#### 3.1.1 Computational Complexity

First, the space complexities of our two methods are \( O(τd) \) caused by the maintenance of \( \mathbf{S}_t, \mathbf{U}, \mathbf{V} \). And the time complexities of our two methods are \( O(τ^2d) \) caused by step 5 in each algorithm which contains the SVD. Thus, both of them are linear in the dimensionality \( d \). Second, compared with the update rules of ADA-LR (3) and RADAGRAD (4), our update rules (7) and (8) do not need random projection.
and are more simple. Third, when \( \varphi \neq 0 \), the computational cost of \( H^{-1}_t \) can still be reduced dramatically, though the updating of \( \beta_t \) may not have a closed-form solution.

### 3.1.2 Theoretical Guarantees

Compared with RADAGRAD, a significant advantage of our methods is that they are equipped with formal theoretical guarantees similar to ADA-FULL. We present regret bounds for Algorithms 1 and 2 in the following two theorems, respectively.

**Theorem 1.** Adaptive dual averaging via frequent directions ensures

\[
R(T) \leq \frac{\delta}{2\eta} \|\beta^*\|_2^2 + \frac{1}{2\eta} \max_{t \leq T} \text{tr}(G_t^{1/2}) \|\beta^*\|_2^2 + \eta \max_{t \leq T} \left( \frac{\max_{t \leq T} \|g_t\|_2 + \sqrt{\Delta_T}}{\delta} \right) \text{tr}(G_t^{1/2}) + \frac{1}{\eta} \sum_{t=1}^T \sqrt{\sigma_t} \max_{t \leq T} \|\beta_{t+1}\|_2^2
\]

where \( \Delta_T = \sum_{t=1}^T \sigma_t \).

**Theorem 2.** Adaptive mirror descent via frequent directions ensures

\[
R(T) \leq \frac{\delta}{2\eta} \|\beta^*\|_2^2 + \frac{1}{2\eta} \max_{t \leq T} \|\beta^* - \beta_t\|_2^2 \text{tr}(G_t^{1/2}) + \eta \max_{t \leq T} \left( \frac{\sqrt{\Delta_T}}{\delta} \right) \text{tr}(G_t^{1/2}) + \frac{\tau}{2\eta} \sum_{t=1}^T \sqrt{\sigma_t} \max_{t \leq T} \|\beta^* - \beta_t\|_2^2
\]

where \( \Delta_T = \sum_{t=1}^T \sigma_t \).

For comparisons, let us introduce the theoretical guarantees of ADA-FULL as follows.

**Theorem 3.** (Theorem 7 of Duchi et al. [26]) Provided with \( \delta \geq \max_{t \leq T} \|g_t\|_2 \), ADA-FULL incorporated into the primal-dual subgradient method ensures

\[
R(T) \leq \frac{\delta}{2\eta} \|\beta^*\|_2^2 + \frac{1}{2\eta} \max_{t \leq T} \|\beta^* - \beta_t\|_2^2 \text{tr}(G_t^{1/2}) + \eta \text{tr}(G_t^{1/2}).
\]

ADA-FULL incorporated into the composite mirror descent method ensures

\[
R(T) \leq \frac{\sigma}{2\eta} \|\beta^*\|_2^2 + \frac{1}{2\eta} \max_{t \leq T} \|\beta^* - \beta_t\|_2^2 \text{tr}(G_t^{1/2}) + \eta \text{tr}(G_t^{1/2}).
\]

Compared with the Theorem 3, we confirm that our methods enjoy the regret bound close to that of ADA-FULL, and we point out the slight differences as following. First, the regret bounds of our two methods contain an additional term, respectively \( \frac{\max_{t \leq T} \|\beta_{t+1}\|_2^2}{2\eta} \) and \( \frac{\max_{t \leq T} \|\beta^* - \beta_t\|_2^2}{2\eta} \). According to Cauchy-Schwarz inequality and \( \Delta_T = \sum_{i=1}^T \sigma_i \), we have

\[
\sum_{t=1}^T \sqrt{\sigma_t} \leq \sqrt{T} \sum_{t=1}^T \sigma_t = \sqrt{T}\Delta_T
\]

which means the additional term does not affect the order of the regret bounds when \( \Delta_T \) is small. Second, our two bounds magnify the third term by \( \max \left( \frac{\max_{t \leq T} \|g_t\|_2 + \sqrt{\Delta_T}}{\delta} \right) \) and \( \max \left( \frac{\sqrt{\Delta_T}}{\delta} \right) \) respectively, which does not change the order of the regret bounds when \( \Delta_T \) is small. Let \( C_k^p \) denote the minimizer of \( \|C_t - C_k^p\|_F \) over all rank \( k \) matrices. According to the property of frequent directions, we have

\[
\Delta_T \leq \|\beta^* - C_k^p\|_2^2 / (\tau - k)
\]

for any \( k < \tau \), which means \( \Delta_T \) is small when the outer product matrix of gradients is approximately low-rank. Moreover, when \( G_T \) is low-rank and \( \text{rank}(G_T) = r \), we have \( \Delta_T = 0 \) by choosing \( \tau = r + 1 \).

Additionally, according to Proposition 2 of Krummenacker et al. [27], the regret bound of ADA-LR is also close to that of ADA-FULL. However, it only holds with high probability and requires \( \tau \geq O(\log d) \). In contrast, the regret bounds of our methods are deterministic and hold for any integer \( \tau > 0 \).

### 3.2 Adaptive Online Learning via FFD

Although the time complexity of ADA-FD is linear in the dimensionality \( d \), it still suffers the quadratic dependence on the sketching size \( \tau \). Motivated by the fast implementation of frequent directions (5), we further accelerate our ADA-FD by doubling the sketching size and reducing the number of time-consuming computations. However, fast frequent directions cannot be directly applied to our methods, because it only computes the SVD of the sketch matrix \( S_t \) once every \( \tau + 1 \) iterations while for updating the decision, our methods need to compute the square root of \( S_t \) and the inverse of \( H_t \) by the SVD of \( S_t \) in each iteration. As a result, besides the doubling trick used in fast frequent directions, we carefully design a more efficient strategy to compute the
Now, we can directly utilize the above two matrices $S_t$ as explained below, which is motivated by the idea of incremental SVD [43].

According to the doubling trick, the size of $S_t$ is changed from $\tau \times d$ to $2\tau \times d$. Moreover, instead of directly maintaining $S_t$, we maintain two matrices $M_t$ and $V_t$ which satisfy

$$M_t^{1/2}V_t^T = S_t \in \mathbb{R}^{2\tau \times d}$$

where $M_t \in \mathbb{R}^{2\tau \times 2\tau}$ is a symmetric matrix and $V_t = [v_1^T, v_2^T, \cdots, v_d^T] \in \mathbb{R}^{d \times 2\tau}$ consists of $r_t \leq 2\tau$ orthonormal vectors and $2\tau - r_t$ zero vectors. For $t = 0$, we simply set $r_0 = 0$, $V_0 = 0_{d \times 2\tau}$, $M_0 = 0_{2\tau \times 2\tau}$.

In each iteration $t = 1, \cdots, T$, after receiving the gradient $g_t$, we first compute

$$g' = V_{t-1}(V_{t-1}^T g_t).$$

If $g' \neq g_t$, we set $r_{t-1} = r_t + 1$ and

$$v_{r_{t-1}}' = \frac{g_t - g'}{\|g_t - g'\|_2}$$

where $(v_i')^T v_i = 0$ for any $1 \leq i < r_{t-1}$. Then, we have $V_{t-1}V_{t-1}^T g_t = g_t$, which leads to

$$V_{t-1}M_{t-1}V_{t-1}^T + g_t g_t^T = V_{t-1}M_{t-1}V_{t-1}^T + V_{t-1}V_{t-1}^T g_t g_t^T V_{t-1}V_{t-1}^T = V_{t-1}M_{t-1}V_{t-1}^T + g_t g_t^T V_{t-1}V_{t-1}^T.$$

So, we can set $r_t = r_{t-1}$, $V_t = V_{t-1}$ and compute $M_t$ as

$$M_t = M_{t-1} + (V_{t-1}^T g_t)(V_{t-1}^T g_t)^T.$$

Now, we can directly utilize the above two matrices $V_t$ and $M_t$ to update the decision. Specifically, we first efficiently compute the SVD of $M_t$ as

$$U \Sigma U^T = M_t.$$

Then, we have

$$S_t = M_t^{1/2}V_t^{1/2} = \sqrt{V_t} \Sigma^{1/2} V_t = V_t^{1/2} U^{1/2} \Sigma U^T V_t^{1/2}$$

which means that $\beta_{t+1}$ can be computed efficiently with Woodbury formula as in ADA-FD. Moreover, according to fast frequent directions, if $r_t = 2\tau$, we need to make the last $\tau + 1$ rows of $S_t$ equal to zero by computing

$$M_t = \max(\Sigma - \sigma I_{2\tau}, 0), V_t = V_t U$$

where $\sigma_t = \Sigma_{\tau\tau}$, and setting $r_t = \tau - 1$ and the right $\tau + 1$ columns of $V_t$ to be zero.

With the above procedures, we propose our second efficient variant of ADA-FULL, named as adaptive online learning via fast frequent directions (ADA-FFD). Furthermore, the detailed procedures of ADA-FFD for the primal-dual subgradient method and composite mirror descent method are summarized in Algorithms 3 and 4, respectively.

### 3.2.1 Computational Complexity

In each iteration, our Algorithms 3 and 4 need to compute $g_t', v_{t-1}'$, $M_t$, the SVD of $M_t$, and $\beta_{t+1}$, the time complexities of which are $O(\tau d)$, $O(d)$, $O(\tau^2)$, and $O(\tau d)$, respectively. Assuming $\tau \leq \sqrt{d}$, in each iteration, the total time complexity of these computations is $O(\tau d)$. Besides these computations, Algorithms 3 and 4 need to compute $V_t = V_t U$ once every $\tau + 1$ iterations, which causes the time complexity of $O(\tau^2 d)$. Therefore, in each iteration, the average time complexities of Algorithms 3 and 4 are

$$O \left( \frac{\tau d T + \tau^2 d (T/\tau)}{T} \right) = O(\tau d)$$

which is linear in both the dimensionality $d$ and the sketching size $\tau$. Meanwhile, the space complexities of Algorithms 3 and 4 are still $O(\tau d)$ caused by the maintenance of $V_t$, because ADA-FDD only doubles the sketching size of ADA-FD. Compared with ADA-FD and RADAGRAD, our ADA-FFD further reduces the time complexity to $O(\tau d)$ while enjoys the same space complexity of $O(\tau d)$.

### 3.2.2 Theoretical Guarantees

Note that Algorithms 3 and 4 enjoy almost the same regret bounds as Algorithms 1 and 2, respectively. For completeness, we present the regret bounds of Algorithms 3 and 4 in the following two theorems, respectively.
Theorem 4. Adaptive dual averaging via fast frequent directions ensures
\[
R(T) \leq \frac{\delta}{2\eta} \|\beta^*\|^2 + \frac{1}{2\eta} \max_{t \leq T} \|g_t\|^2 \|\beta^*\|^2 + \eta \max_{t \leq T} \left( \frac{\sqrt{\Delta_t}}{\delta} \right) \text{tr}(G_T^{1/2}) + \eta \max_{t \leq T} \|\beta_{t+1}\|^2
\]
where \(\Delta_t = \sum_{i=1}^T \sigma_i\).

Theorem 5. Adaptive mirror descent via fast frequent directions ensures
\[
R(T) \leq \frac{\delta}{2\eta} \|\beta^*\|^2 + \frac{1}{2\eta} \max_{t \leq T} \|\beta^* - \beta_t\|^2 \text{tr}(G_T^{1/2}) + \eta \max_{t \leq T} \left( \frac{\sqrt{\Delta_t}}{\delta} \right) \text{tr}(G_T^{1/2}) + \frac{\tau}{\eta} \sum_{t=1}^T \sqrt{\sigma_t} \max_{t \leq T} \|\beta^* - \beta_t\|^2
\]
where \(\Delta_t = \sum_{i=1}^T \sigma_i\).

According to the property of fast frequent directions proved by Ghashami et al. [31], we still have
\[
\Delta_T \leq \|C_T - C_T^*\|^2/(\tau - k)
\]
for any \(k < \tau\). At first sight, (11) is the same as (9) in Theorem 1 and (12) only doubles the last term of (10) in Theorem 2. However, the sketching size of ADA-FFD is \(2\tau\). For a fair comparison, we should reduce it to \(\tau\), which implies that \(\tau\) in (12) and (13) would be replaced with \(\tau/2\).

3.3 Discussions

We would like to emphasize the differences between our work and Krummenacher et al. [27]. First, both ADA-LR and RADAGRAD only consider the composite mirror descent method, ignoring the primal-dual subgradient method. Second, unlike random projections adopted in Krummenacher et al. [27], our work makes use of frequent directions and fast frequent directions that belong to deterministic matrix sketching techniques. As a result, the theoretical guarantees of our methods are deterministic which are more robust than the probabilistic regret bound of ADA-LR. Third, our methods are very efficient in the sense that the time and storage complexities are linear in the dimensionality \(d\), without losing the regret bound. Especially, the time complexities of our Algorithms 3 and 4 are also linear in the sketching size \(\tau\).

By contrast, RADAGRAD only reaches similar complexities as our Algorithms 1 and 2 while does not have theoretical justifications.

Next, we discuss the differences between our work and Luo et al. [30]. First, our methods and Luo et al. [30] are designed for different tasks. Our goal is to develop an efficient variant of ADA-FULL, and the goal of Luo et al. [30] is to accelerate ONS. Note that ADA-FULL is a data dependent algorithm for general convex functions and ONS is proposed to derive a logarithmic regret bound for exponentially concave functions. Second, the theoretical analysis in our work is obviously different from Luo et al. [30]. They provide a regret bound of ONS with order \(O(\sqrt{T})\) for general convex functions by setting \(\eta_t = O(1/\sqrt{T})\). Although its update rules (6) can be reformulated as
\[
H_t = \delta I_d + \sum_{i=1}^T \frac{1}{\sqrt{d}} g_i g_i^T \quad \text{and} \quad \beta_{t+1} = \beta_t - \eta H_t^{-1} g_t
\]
which is similar to ADA-FULL, its \(O(\sqrt{T})\) bound destroys the data dependent property. Third, the main algorithm of Luo et al. [30] can be regarded as a mirror descent method. In contrast, our work proposes methods for both the primal-dual subgradient method and composite mirror descent method, respectively.

4 THEORETICAL ANALYSIS

In this section, we provide the proofs of Theorems 1 and 2 and omit the proofs of Theorems 4 and 5, because it is not difficult to prove Theorems 4 and 5 in the same way.

4.1 Supporting Results

The following results are used throughout our analysis.

Lemma 1. (Variant of Proposition 2 in Duchi et al. [26].) Let the sequence \(\{\beta_t\}\) be generated by Algorithm 1. We have
\[
R(T) \leq \frac{1}{\eta} \Psi_T(\beta^*) + \frac{\eta}{2} \sum_{t=1}^T \|f'_t(\beta_t)\|^2_{\psi_{t-1}} + \frac{\sum_{t=1}^T \sqrt{\sigma_t}}{2\eta} \max_{t \leq T} \|\beta_{t+1}\|^2
\]

Lemma 1 can be regarded as a variant of Proposition 2 in Duchi et al. [26], when the condition \(\Psi_{t+1}(\beta) \geq \Psi_t(\beta)\) cannot be met due to \(H_{t+1} \neq H_t\) in this work. And it can be derived from the proof of Proposition 2 in Duchi et al. [26] with slight modifications to deal with \(\Psi_{t+1}(\beta) \geq \Psi_t(\beta)\). We include the proof for completeness.

Proof. The conjugate dual of \(t\phi(\beta) + \frac{1}{\eta} \Psi_t(\beta)\) is defined by
\[
\Phi_t^*(g) = \sup_{\beta} \left\{ \langle g, \beta \rangle - t\phi(\beta) - \frac{1}{\eta} \Psi_t(\beta) \right\}
\]
Thus, the gradient of \(\Phi_t^*(g)\) can be calculated as
\[
\nabla \Phi_t^*(g) = \arg \min_{\beta} \left\{ -\langle g, \beta \rangle + t\phi(\beta) + \frac{1}{\eta} \Psi_t(\beta) \right\}
\]
Because \(\frac{1}{\eta} \Psi_t(\beta)\) is \(\frac{1}{\eta}\)-strongly convex with respect to the norm \(\|\cdot\|_{\psi_t}\), we have \(\|\nabla \Phi_t^*(x) - \nabla \Phi_t^*(y)\|_{\psi_t} \leq \eta\|x - y\|_{\psi_t}\) which means the function \(\Phi_t^*\) has \(\eta\)-Lipschitz continuous gradients with respect to \(\|\cdot\|_{\psi_t}\). Further, we have
\[
\Phi_t^*(y) \leq \Phi_t^*(x) + \langle \nabla \Phi_t^*(x), y - x \rangle + \frac{\eta}{2} \|y - x\|_{\psi_t}^2
\]
Both the identity (14) and the bound (15) were used in the proof of Proposition 2 in Duchi et al. [26]. In order to complete the proof, we introduce an inequality
\[
\sum_{t=1}^T \{f_t(\beta_t) + \phi(\beta_t) - f_t(\beta^*) - \phi(\beta^*)\} \leq \frac{1}{\eta} \Psi_T(\beta^*) + \sum_{t=1}^T \{\langle g_t, \beta_t \rangle + \phi(\beta_t)\} + \Phi_T^*(-\bar{g}_T)
\]
from the proof of Proposition 2 in Duchi et al. [26] again.

Let \( \Sigma, \Sigma_i, V_i \) denote the values of \( \Sigma, \Sigma_i, V \) in the \( t \)-th round of Algorithm 1. For each \( t = 1, \ldots, T \), we have

\[
(S_t^r S_t)^{1/2} - (S_{t-1}^r S_{t-1})^{1/2} + \sqrt{\sigma_t V_t V_t^T} - \sqrt{(S_{t-1}^r S_{t-1})^{1/2}} \\
\geq \sqrt{V_t V_t^T} - (S_{t-1}^r S_{t-1})^{1/2} \\
= (S_{t-1}^r S_{t-1} + g_t g_t^T)^{1/2} - (S_{t-1}^r S_{t-1})^{1/2} \\
\geq 0
\]

which implies \(-\Psi_t(x) \leq -\Psi_{t-1}(x) + \sqrt{\sigma_t} \|x\|^2_2\).

Thus, we have

\[
\Phi_T^*(-g_T) = -(g_T, \beta_{T+1}) - T\varphi(\beta_{T+1}) - \frac{1}{\eta} \Psi_T(\beta_{T+1}) \\
\leq -(g_T, \beta_{T+1}) - T\varphi(\beta_{T+1}) - \frac{1}{\eta} \Psi_{t-1}(\beta_{T+1}) \\
+ \frac{\sqrt{\sigma_t}}{2\eta} \|\beta_{T+1}\|^2_2 \\
\leq \sup_\beta \left( -(g_T, \beta) - (T-1)\varphi(\beta) - \frac{1}{\eta} \Psi_{t-1}(\beta) \right) \\
- \varphi(\beta_{T+1}) + \frac{\sqrt{\sigma_t}}{2\eta} \|\beta_{T+1}\|^2_2 \\
= \Phi_{t-1}^*(-g_T) - \varphi(\beta_{T+1}) + \frac{\sqrt{\sigma_t}}{2\eta} \|\beta_{T+1}\|^2_2
\]

which contains an additional term \(\frac{\sqrt{\sigma_t}}{2\eta} \|\beta_{T+1}\|^2_2\) caused by \(H_T \neq H_{T-1}\) comparing with Duchi et al. [26].

Using the identity (14), the bound (15) and the inequality (16), we have

\[
\sum_{t=1}^T \{f_t(\beta_t) + \varphi(\beta_{t+1}) - f_t(\beta^*) - \varphi(\beta^*)\} \\
\leq \frac{1}{\eta} \Psi_T(\beta^*) + \frac{1}{\eta} \sum_{t=1}^T \{\langle g_t, \beta_t \rangle + \varphi(\beta_{t+1})\} \\
+ \frac{1}{\eta} \sum_{t=1}^T \{\langle g_t, \beta_t \rangle + \varphi(\beta_{t+1})\} \\
+ \frac{\sqrt{\sigma_t}}{2\eta} \|\beta_{T+1}\|^2_2 \\
\leq \frac{1}{\eta} \Psi_T(\beta^*) + \frac{1}{\eta} \sum_{t=1}^T \{\langle g_t, \beta_t \rangle + \varphi(\beta_{t+1})\} \\
+ \frac{\sqrt{\sigma_t}}{2\eta} \|\beta_{T+1}\|^2_2 \\
\leq \frac{1}{\eta} \Psi_T(\beta^*) + \frac{1}{\eta} \sum_{t=1}^{T-1} \{\langle g_t, \beta_t \rangle + \varphi(\beta_{t+1})\} \\
+ \frac{\sqrt{\sigma_t}}{2\eta} \|\beta_{T+1}\|^2_2
\]

By repeating the above steps, we have

\[
\sum_{t=1}^T \{f_t(\beta_t) + \varphi(\beta_{t+1}) - f_t(\beta^*) - \varphi(\beta^*)\} \\
\leq \frac{1}{\eta} \Psi_T(\beta^*) + \frac{1}{\eta} \sum_{t=1}^T \|g_t\|_2^2 \varphi(\beta_{t+1}) + \frac{\sqrt{\sigma_t}}{2\eta} \|\beta_{T+1}\|^2_2 \\
+ \Phi_0^*(-g_0).
\]

Note that \(\varphi(\beta) = 0\) and \(\Phi_0^*(0) = 0\). We complete the proof.

Lemma 2. (Proposition 3 in Duchi et al. [26]). Let the sequence \(\{\beta_t\}\) be generated by Algorithm 2. We have

\[
R(T) \leq \frac{1}{\eta} \sum_{t=1}^T \{B_{\Psi_{t+1}}(\beta^*, \beta_{t+1}) - B_{\Psi_t}(\beta^*, \beta_{t+1})\} \\
+ \frac{1}{\eta} B_{\Psi_t}(\beta^*, \beta_1) + \sum_{t=1}^T \|f'_t(\beta_t)\|_2^2.
\]

Lemma 3. (Lemma 10 in Duchi et al. [26]) Let \( G_i = \sum_{i=1}^t g_i g_i^T \) and \( A^1 \) denote the pseudo-inverse of \( A \), then

\[
\sum_{t=1}^T \langle g_t, (G_t^{1/2})^T g_t \rangle \leq 2 \sum_{t=1}^T \langle g_t, (G_t^{1/2})^T g_t \rangle = 2 \text{tr}(G_t^{1/2}).
\]

Lemma 4. (Derived From Theorem 3.1 and Its Proof in Ghoshami et al. [31]) Let \( \Delta_t = \sum_{i=1}^t \sigma_i \). In Algorithms 1 and 2, \( S_t \) is the sketch generated by performing frequent directions on the input \( C_t \). Then for any \( t \) and \( k < T \), \( C_t^T C_t \geq S_t^r S_t \geq C_t^T C_t - \Delta_t I_2 \) and \( \Delta_t \leq \|C_t - C_t^k\|_F^2/\gamma - k \), where \( C_t^k \) denotes the minimizer of \( \|C_t - C_t^k\|_F \) over all rank \( k \) matrices.

4.2 Proof of Theorem 1

We first consider \( \frac{1}{\eta} \Psi_T(\beta^*) \) in the upper bound of Lemma 1. We have

\[
\frac{1}{\eta} \Psi_T(\beta^*) = \frac{1}{2\eta} \langle \beta^*, (\delta I_d + (S_t^r S_t)^{1/2}) \beta^* \rangle \\
\leq \frac{\delta}{2\eta} \|\beta^*\|^2_2 + \frac{1}{2\eta} \langle \beta^*, (C_t^{1/2} C_t^{1/2})\beta^* \rangle \\
\leq \frac{\delta}{2\eta} \|\beta^*\|^2_2 + \frac{1}{2\eta} \lambda_{\text{max}}(C_t^{1/2}) \|\beta^*\|^2_2 \\
\leq \frac{\delta}{2\eta} \|\beta^*\|^2_2 + \frac{1}{2\eta} \text{tr}(G_t^{1/2}) \|\beta^*\|^2_2
\]

where \(\lambda_{\text{max}}(C_t^{1/2})\) denotes the largest eigenvalue of \(C_t^{1/2}\).

Before considering \( \frac{1}{\eta} \sum_{t=1}^T \|f'_t(\beta_t)\|_2^2 \), we need to derive the lower bound of \(H_{t-1}\). Let \( c = \frac{\delta}{\|g_t\|_2 + \sqrt{\Delta_{t-1}}} \). If \( c < 1 \), we have

\[
H_{t-1} = \delta I_d + (S_{t-1}^r S_{t-1})^{1/2} \\
\geq c \|g_t\|_2 I_d + \sqrt{\Delta_{t-1} I_d + (S_{t-1}^r S_{t-1})^{1/2}} \\
\geq c \|g_t\|_2 I_d + (\Delta_{t-1} I_d + S_{t-1}^r S_{t-1})^{1/2} \\
\geq c \|g_t\|_2 I_d + (C_{t-1}^T C_{t-1})^{1/2} \\
\geq c \|g_t\|_2 I_d + (C_{t-1}^T C_{t-1})^{1/2}
\]

where the second inequality is due to \( \sqrt{\Delta_{t-1} + x^2} \geq \sqrt{\Delta_{t-1} + x^2} \) for any \( x \geq 0 \) and the third inequality is due to Lemma 4.

In the other case \( \delta \geq \sqrt{\Delta_{t-1} + \|g_t\|_2^2} \), we have

\[
H_{t-1} = \delta I_d + (S_{t-1}^r S_{t-1})^{1/2} \\
\geq \|g_t\|_2 I_d + \sqrt{\Delta_{t-1} I_d + (S_{t-1}^r S_{t-1})^{1/2}} \\
\geq \|g_t\|_2 I_d + (\Delta_{t-1} I_d + S_{t-1}^r S_{t-1})^{1/2} \\
\geq \|g_t\|_2 I_d + (C_{t-1}^T C_{t-1})^{1/2} \\
\geq (C_{t-1}^T C_{t-1})^{1/2}.
\]
Let $t$ be an integer such that $1 \leq t \leq T$. Then, we have

$$
\sum_{t=1}^{T} \| f_t'(\beta_t) \|^2_{2} = \sum_{t=1}^{T} \langle g_t, (H_{t-1})^{-1} g_t \rangle \\
\leq \sum_{t=1}^{T} \max \left( 1, \frac{\| g_t \|^2 + \sqrt{\Delta t}}{\delta} \right) \langle g_t, (G_t^*)^2 g_t \rangle \\
\leq \max \left( 1, \frac{\max \{ \| g_t \|^2 + \sqrt{\Delta t} \}}{\delta} \right) \sum_{t=1}^{T} \langle g_t, (G_t^*)^2 g_t \rangle \\
\leq 2 \max \left( 1, \frac{\max \{ \| g_t \|^2 + \sqrt{\Delta t} \}}{\delta} \right) \tr(G_T^2) \tag{18}
$$

where the last inequality is due to Lemma 3.

We complete the proof by substituting (17) and (18) into Lemma 1.

### 4.3 Proof of Theorem 2

Let $\Sigma_t, \Sigma'_t, V_t$ denote the values of $\Sigma, \Sigma', V$ in the $t$-th round of Algorithm 2.

For each $t = 1, \cdots, T$, we have

$$
\langle g_t, (H_{t-1})^{-1} g_t \rangle \\
= \langle g_t, H_t^{-1} g_t \rangle - \langle g_t, H_{t-1}^{-1} g_t \rangle \\
\geq \tr(\tilde{G}_t^2) \tag{19}
$$

where the second inequality is due to $\sqrt{\Delta t} + x \geq \sqrt{\Delta t + x^2}$ for $x \geq 0$ and the third inequality is due to Lemma 4.

And in the other case $\delta \geq \sqrt{\Delta t}$, we have

$$
H_t = \delta I_d + (S_t^T S_t)^{1/2} \geq \frac{\delta}{\sqrt{\Delta t}} (C_t^T C_t)^{1/2}
$$

Thus, for any $\delta > 0$, we have

$$
H_t \succeq \min \left( 1, \frac{\delta}{\sqrt{\Delta t}} \right) (C_t^T C_t)^{1/2}. \tag{20}
$$

Then, we have

$$
\sum_{t=1}^{T} \| f_t'(\beta_t) \|^2_{2} = \sum_{t=1}^{T} \langle g_t, (H_{t-1})^{-1} g_t \rangle \\
\leq \sum_{t=1}^{T} \max \left( 1, \frac{\sqrt{\Delta t}}{\delta} \right) \langle g_t, (G_t^*)^2 g_t \rangle \\
\leq \max \left( 1, \frac{\sqrt{\Delta t}}{\delta} \right) \sum_{t=1}^{T} \langle g_t, (G_t^*)^2 g_t \rangle \\
\leq 2 \max \left( 1, \frac{\sqrt{\Delta t}}{\delta} \right) \tr(G_T^2) \tag{20}
$$

where the last inequality is due to Lemma 3.

By substituting (19) and (20) into Lemma 2, we complete the proof.

### 5 Experiments

In this section, we perform numerical experiments on online regression, online classification, training CNN, and fine-tuning much deeper CNN to verify the efficiency and effectiveness of our algorithms.
Fig. 1. The comparison of regret among different algorithms with the default $\tau = 20$ on synthetic data for the PDS method and CMD method.

Fig. 2. The comparison of regret among different algorithms with different $\eta$ on synthetic data for the PDS method and CMD method.

5.1 Baselines and the Default Setup

First, we show that ADA-FD and ADA-FFD approximate ADA-FULL well, and outperform ADA-DIAG. Second, we compare our methods with existing approximations of ADA-FULL including ADA-LR and RADAGRAD. Third, we compare our methods with FD-SON [30] that is an approximation of ONS by frequent directions. Fourth, we compare our algorithms with the classical online gradient descent (OGD) [44] which updates as $\beta_{t+1} = \beta_t - \frac{\eta}{T} \nabla f_t$ and ADAM [45] that is a popular method for training neural networks. Note that ADAM have four parameters including $\eta$, a constant $\delta$, and two decay rates $\beta_1, \beta_2$.

In all experiments, by default, we search the step size $\eta$ for all algorithms from the set of $\{1e^{-4}, 1e^{-3}, \cdots, 1\}$, and choose the value that achieves the best regret for experiments on online regression or the best training performance for other experiments. To simplify the parameter settings and control the computational cost, we simply set $\tau = 20$ for algorithms using matrix approximation by default. Moreover, we set $\delta = 1e^{-8}$, $\beta_1 = 0.9$, $\beta_2 = 0.999$ for ADAM as recommended by Kingma and Ba [45], and also set $\delta = 1e^{-8}$ for ADA-DIAG. For ADA-FULL, ADA-LR, RADAGRAD, and FD-SON, we select the parameter $\delta$ from the set of $\{1e^{-4}, 1e^{-3}, \cdots, 10\}$, and choose the value that achieves the best regret for experiments on online regression or the best training performance for other experiments. For our ADA-FD and ADA-FFD with the default $\tau = 20$, we set $\delta = 1$ when combining with the composite mirror descent (CMD) method, and set $\delta = 10$ when combining with the primal-dual subgradient (PDS) method.

In the above way, by default, our ADA-FD and ADA-FFD only have one varying hyperparameter—the step size $\eta$ in different experiments.

5.2 Online Regression

We first consider the problem of online regression where $f_t(\beta) = \|\beta^T x_t - y_t\|$ and $y_t = \beta^T x_t$ with ideal synthetic data. Specifically, we set $T = 10,000$, $d = 500$, and $\beta_t = \beta_0/\|\beta_0\|_2$, where each entry of $\beta_0$ is drawn independently from $N(0, 1)$. In order to meet the requirement of the approximately low-rank structure, each data point $x_t$ is sampled independently from a Gaussian distribution $N(\mu, \Sigma)$ where $\mu = 1$ and $\Sigma$ has rapidly decaying eigenvalues $\lambda_j(\Sigma) = \lambda_0 10^{-j\alpha}$ with $\alpha = 2$ and $\lambda_0 = 100$. For the PDS method, we compare ADA-FD and ADA-FFD against ADA-FULL and ADA-DIAG. And for the CMD method, we compare ADA-FD and ADA-FFD against all baselines. For fairness, all algorithms are run with the same permutation of the function sequence. Due to the randomness of ADA-LR and RADAGRAD, their experiments are run 20 times, and we report the average regret with standard deviations.

Fig. 1 shows the comparison of regret among different algorithms with the default $\tau = 20$. For both the PDS method and CMD method, our two algorithms outperform ADA-DIAG, and are close to ADA-FULL. For the CMD method, our two algorithms outperform ADA-LR, RADAGRAD, FD-SON, OGD, and ADAM. Note that according to the default setup, we need to search the best step size $\eta$ from the set of $\{1e^{-4}, 1e^{-3}, \cdots, 1\}$ for all algorithms. So, it is reasonable to investigate whether our algorithms and other baselines are similarly sensitive to the value of $\eta$. To this end, we further present the comparison of the total regret among different algorithms with different $\eta$ in Fig. 2. We find that all algorithms are sensitive to the value of $\eta$. So, it is fair to search the best value of $\eta$ for each algorithm.

Fig. 3 shows the total regret and running time of all algorithms with $\tau = \{20, 40, 60, 80\}$ for the CMD method.
This article has been accepted for publication in a future issue of this journal, but has not been fully edited. Content may change prior to final publication. Citation information: DOI 10.1109/TPAMI.2021.3096880, IEEE Transactions on Pattern Analysis and Machine Intelligence.

Fig. 5. The comparison of mistakes and test accuracy among different algorithms on Gisette for the PDS method and CMD method.

Fig. 6. The comparison of mistakes and test accuracy among different algorithms on Epsilon for the PDS method and CMD method.

Note that the parameter $\delta$ of our two algorithms with $\tau > 20$ is selected according to the same way used for ADA-FULL. From this figure, we have the following observations:

- When $\tau$ is increased from 20 to 80, the total regret of algorithms using matrix approximation is decreased, and our ADA-FD and ADA-FFD always outperform ADA-LR, RADAGRAD, and FD-SON. The running time of each algorithm using matrix approximation increases as $\tau$ becomes larger. ADA-LR is much slower than all other methods except ADA-FULL. Our ADA-FD is significantly faster than ADA-FULL and as fast as RADAGRAD and FD-SON. Moreover, ADA-FFD is faster than ADA-FD.

- With $\tau = 80$, our ADA-FD and ADA-FFD can slightly outperform ADA-FULL. It is reasonable, because our ADA-FD and ADA-FFD maintain a strictly low-rank approximation of gradient outer products, which could eliminate possible noises in the approximately low-rank data.

Furthermore, it is obvious that our ADA-FD with $\tau = 501$ and ADA-FFD with $\tau = 251$ will have the same regret as ADA-FULL. Therefore, larger $\tau$ might not be better for minimizing the regret, which is clearly verified by Fig. 4. Moreover, from Fig. 4, we notice that the regret of ADA-FFD suddenly degrades back to that of ADA-FULL when $\tau$ changes from 250 to 251. It is reasonable, because the process of ADA-FFD with $\tau = 251$ significantly differs from that of ADA-FFD with $\tau = 250$. For ADA-FFD with $\tau = 251$, when $t$ increases, the rank of $V_t M_t V_t^\top$ first gradually increases from 1 to 500, and then is always kept at 500. However, with $\tau = 250$, every time the rank of $V_t M_t V_t^\top$ reaches $2\tau = 500$, the $if$ statement in step 12 of Algorithm 4 is triggered. In such cases, the rank of $V_t M_t V_t^\top$ will be reduced to $\tau - 1 = 249$, because $\tau + 1 = 251$ columns of the matrix $V_t M_t^{1/2}$ are zeroed out. So, for ADA-FFD with $\tau = 250$, there exist cyclical changes about the rank of $V_t M_t V_t^\top$: it repeatedly grows from $\tau - 1$ to $2\tau$.

Since the regret of ADA-FD is close to that of ADA-FFD, to facilitate presentations, we only report the results of ADA-FFD in the following.

5.3 Online Classification

Then, we perform online classification to evaluate the performance of our ADA-FFD with two real world datasets from LIBSVM repository [46]: Gisette and Epsilon which are high-dimensional and dense. At each round, the learning algorithm receives a single example $(x_t, y_t)$ and suffers the squared hinge loss $f_t(\beta) = (\max(0, 1 - y_t \beta^\top x_t))^2$. Each algorithm ends with a single pass through the training data. Following Duchi et al. [26], we adopt two metrics to measure the performance: the online mistakes and the offline accuracy on the testing data. For the PDS method, we compare ADA-FFD against ADA-DIAG. And for the CMD method, we compare ADA-FFD against ADA-DIAG, RADAGRAD, FD-SON, OGD, and ADAM. We omit the results of ADA-FULL and ADA-LR, because they are too slow. Both datasets are divided into the training part and testing part, and the numbers of training and testing examples are shown in Table 1.

For both datasets, we repeat the experiments of all algorithms 20 times with random permutations of training data, and report the average mistakes and test accuracy with standard deviations. Fig. 5 and 6 show the comparison of mistakes and test accuracy among different algorithms. We find that our ADA-FFD outperforms ADA-DIAG, RADAGRAD, FD-SON, OGD, and ADAM in terms of mistakes on Gisette and test accuracy on both datasets. In terms
of mistakes on Epsilon, our ADA-FFD outperforms ADA-DIAG, RADAGRAD, OGD, and ADAM. Fig. 7 shows the comparison of running time among different algorithms.

Our ADA-FFD is faster than RADAGRAD and FD-SON when $d = 5000$ and $d = 2000$. Compared with Fig. 3(b), the advantage of our ADA-FFD in running time is more obvious for the larger $d$.

### 5.4 Non-convex Optimization in CNN

Besides convex optimization, ADA-DIAG incorporated into the CMD method has also been applied to non-convex optimization such as training neural networks and has performed well. Therefore, we take convolutional neural networks as an example and compare ADA-FFD against ADA-DIAG, RADAGRAD, FD-SON, OGD, and ADAM on training classical neural networks. We consider a 4-layer CNN from the Keras examples directory, which contains two $3 \times 3$ convolutional layers and two fully connected layers. The first convolutional layer generates 32 channels, and the second convolutional layer generates 64 channels followed by $2 \times 2$ max-pooling and 0.25 dropout. The first fully connected layer contains 128 hidden units followed by 0.5 dropout, and the second fully connected layer is a logistic layer. The activation functions of all layers except the logistic layer are ReLU, and the objective is the cross-entropy loss. We use this simple and standard architecture.
to perform classification on the MNIST [47], CIFAR10 [48], CIFAR100 [48], and SVHN [49] datasets.

Following Krummenacher et al. [27], we only consider applying ADA-FFD, RADAGRAD, and FD-SON to the convolutional layers, and other layers are still trained with ADA-DIAG. For all algorithms, we run 20 times with the batch size 128, and report the average training loss and test accuracy with standard deviations. Fig. 8 shows the comparison of training loss and test accuracy during training among different algorithms. We find that our ADA-FFD outperforms ADA-DIAG, RADAGRAD, FD-SON, and OGD on all datasets when applied to training CNN. Moreover, our ADA-FFD outperforms ADAM on CIFAR 10 and CIFAR100, and is close to ADAM on MNIST and SVHN. We also note that RADAGRAD is outperformed by ADA-FFD when the same running time is used, which validates that our ADA-FFD has a better ability to approximate ADA-FULL than RADAGRAD. To make a clear comparison about the time complexity, Fig. 9 shows the comparison of running time costed by one epoch of each algorithm. We find that even if all algorithms run with the same number of epochs, our ADA-FFD is only a little slower than ADA-DIAG, OGD, and ADAM.

5.5 Fine-tuning Much Deeper CNN

We also consider the problem of fine-tuning VGG19 [50] to adapt for CIFAR10, which contains sixteen convolutional layers and three fully connected layers. Note that VGG19 is originally designed for ImageNet [51] with 1,000 classes, while CIFAR10 only has 10 classes. Therefore, we replace the three fully connected layers of VGG19 with one fully connected layer that contains 10 hidden units and adopts softmax as the activation function. Furthermore, because $32 \times 32$ images in CIFAR10 are invalid inputs for VGG19, we resize all images in CIFAR10 to $64 \times 64$. Before fine-tuning the modified VGG19, we initialize its convolutional layers with the weights of VGG19 pre-trained on ImageNet. Then, we freeze the weights of all convolutional layers, and only train the top fully connected layer using the resized CIFAR10.

For all algorithms, we run 20 times with the batch size 128, and report the average training loss and test accuracy with standard deviations. Fig. 10 shows the comparison of training loss, test accuracy, and running time of one epoch among different algorithms on fine-tuning VGG19 for CIFAR10. In the aspects of training loss and test accuracy, we find that our ADA-FFD outperforms ADA-DIAG, RADAGRAD, and OGD, and is slightly better than FD-SON. Although the training loss of our ADA-FFD is worse than that of ADAM, the test accuracy of ADA-FFD is better. Moreover, if all algorithms run with the same number of epochs, our ADA-FFD is only a little slower than ADA-DIAG, OGD, and ADAM.

6 Conclusions

In this paper, we first present ADA-FD to approximate ADA-FULL using frequent directions, which reduces the time and space complexities to $O(\tau^2 d)$ and $O(\tau d)$, respectively, and has the similar efficiency compared to ADA-DIAG. Then, we further present ADA-FFD to reduce the average time complexity to $O(\tau d)$ that is linear in both the dimensionality $d$ and the sketching size $\tau$, which only needs to double the space complexity of ADA-FD for accelerating frequent directions used in it. Furthermore, according to our theoretical analysis, our ADA-FD and ADA-FFD enjoy the regret bound close to that of ADA-FULL when the outer product matrix of gradients is approximately low-rank. Numerical experiments on different problems and datasets demonstrate the efficiency and effectiveness of our algorithms.

Acknowledgments

This work was partially supported by the NSFC (61976112), NSFC-NRF Joint Research Project (61861146001), and JiangsuSF (BK20200064).

References


Yuanyu Wan received the B.S. degree in Software Engineering from Sun Yat-Sen University, China, in 2016 and the M.S. degree in Computer Science and Technology from Nanjing University, China, in 2019. He is currently working toward the Ph.D. degree in the Department of Computer Science and Technology, Nanjing University, China. His research interests include machine learning and optimization.

Lijun Zhang received the B.S. and Ph.D. degrees in Software Engineering and Computer Science from Zhejiang University, China, in 2007 and 2012, respectively. He is currently a Research Professor of the School of Artificial Intelligence, Nanjing University, China. Prior to joining Nanjing University, he was a postdoctoral researcher at the Department of Computer Science and Engineering, Michigan State University, USA. His research interests include machine learning, optimization, information retrieval and data mining.